## Statistics of distances in planar maps

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maps and distances: generalities





$$
\text { (D) } D
$$


$\diamond$ vertices (here of degree 4)

$\diamond$ vertices
$\diamond$ edges (with possibly loops or multiple edges)

$\diamond$ vertices
$\diamond$ edges
$\diamond$ faces with a single boundary (here of degree 3)

$\diamond$ vertices
$\diamond$ edges
$\diamond$ faces
$\rightarrow$ pointed maps

$\diamond$ vertices
$\diamond$ edges
$\diamond$ faces
$\rightarrow$ rooted maps
dual (rooted) map

dual (rooted) map

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coding of a (rooted) map by a (rooted) quadrangulation

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$$




## simple enumeration problems

enumerate, say planar quadrangulations with $F$ faces


## distance statistics

enumerate, say planar quadrangulations with $F$ faces

and with 2 marked vertices

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enumerate, say planar quadrangulations with $F$ faces

and with 2 marked vertices

## distance statistics

enumerate, say planar quadrangulations with $F$ faces

and with 2 marked vertices at prescribed distance
$\rightarrow$ distance profile


enumerate, say planar quadrangulations with $F$ faces

and with 3 marked vertices
enumerate, say planar quadrangulations with $F$ faces

and with 3 marked vertices
enumerate, say planar quadrangulations with $F$ faces

and with 3 marked vertices at prescribed pairwise distances

## number of geodesics

enumerate, say planar quadrangulations with $F$ faces

with 2 marked vertices

## number of geodesics

enumerate, say planar quadrangulations with $F$ faces

with 2 marked vertices
with marked geodesic paths $\rightarrow$ number of geodesic paths
$\log \frac{\langle \# \text { geods }\rangle_{d}}{\langle \# \text { points }\rangle_{d}}$


## $\log \frac{\langle \# \text { geods }\rangle_{d}}{\langle \# \text { points }\rangle_{d}}$


the bijection with mobiles

## from maps to well-labeled mobiles

starting from a pointed planar map with even-valent faces


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## from maps to well-labeled mobiles

starting from a pointed planar map with even-valent faces

end up with a well-labeled mobile

## well-labeled mobiles

## well-labeled:



## well-labeled mobiles

## well-labeled:

(i) positive integer labels


## well-labeled mobiles

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(i) positive integer labels
(ii) at least one label 1


## well-labeled mobiles

## well-labeled:

(i) positive integer labels
(ii) at least one label 1
(iii) rules on labels


## well-labeled mobiles $\rightarrow$ maps

going clockwise around the tree, each corner $\ell$ has a successor $\ell$ - 1


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# other species of trees mobilaceae family 

## arbitrary degrees

start with a pointed planar map


## arbitrary degrees

start with a pointed planar map


## arbitrary degrees

start with a pointed planar map


## arbitrary degrees

start with a pointed planar map
(0)

end up with a new type of mobile



## eulerian maps

start with an eulerian (face bi-colored) planar map


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end up with a new type of mobile


## eulerian maps with blocked edges

start with an eulerian planar map with blocked edges


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end up with a new type of mobile

## eulerian maps with hard particles

$\diamond$ Consider eulerian maps with at most 1 particle per face $\diamond$ Decide to block or not edges between two occupied faces

$\diamond$ Weight -1 per blocked edge $\rightarrow$ selects hard-particle configurations

## generating functions for quadrangulations

## case of quadrangulations



Schaeffer's bijection

## quadrangulations $\rightarrow$ well-labeled trees

start with a pointed planar quadrangulation


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## quadrangulations $\rightarrow$ well-labeled trees

start with a pointed planar quadrangulation
(0)


## quadrangulations $\rightarrow$ well-labeled trees

start with a pointed planar quadrangulation

end up with a planar well-labeled tree

## well-labeled trees

well-labeled:


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## well-labeled:

(i) positive integer labels


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(ii) there is at least one label 1


## well-labeled trees

## well-labeled:

(i) positive integer labels
(ii) there is at least one label 1
(iii) labels vary by at most 1 between neighbors


## well-labeled trees $\rightarrow$ quadrangulations

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## map-tree correspondence

pointed planar quadrangulation
well-labeled tree
(with an origin vertex)

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pointed planar quadrangulation
(with an origin vertex)
vertices at distance $\ell$ from the origin
well-labeled tree
vertices labeled $\ell$

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pointed planar quadrangulation (with an origin vertex)
vertices at distance $\ell$ from the origin
edges $(\ell-1) \leftrightarrow \ell$

well-labeled tree
vertices labeled $\ell$
corner labeled $\ell$

## map-tree correspondence

pointed planar quadrangulation
(with an origin vertex)
vertices at distance $\ell$ from the origin
edges $(\ell-1) \leftrightarrow \ell$
well-labeled tree
vertices labeled $\ell$
corner labeled $\ell$
planted at a corner labeled $\ell$

## map-tree correspondence

pointed planar quadrangulation
(with an origin vertex)
vertices at distance $\ell$ from the origin
edges $(\ell-1) \leftrightarrow \ell$
corner labeled $\ell$
well-labeled tree
vertices labeled $\ell$
marked edge $(\ell-1) \leftrightarrow \ell \quad$ planted at a corner labeled $\ell$
rooted planar quadrangulation (with a root edge)

## generating functions

## well-labeled:

(i) positive integer labels
(ii) there is at least one label 1
(iii) labels vary by at most 1 between neighbors

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## generating functions

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- with cond. (ii) $\quad \rightarrow G_{\ell}(g)=R_{\ell}(g)-R_{\ell-1}(g), \quad R_{0} \equiv 0$
$\rightarrow G_{1}=R_{1}$ : gen. func. for rooted planar quadrangulations


## recursion relations

$$
R_{\ell}=\frac{1}{1-g\left(R_{\ell+1}+R_{\ell}+R_{\ell-1}\right)}
$$

with $R_{0}=0$.
$R_{\ell} \xrightarrow{\ell \rightarrow \infty} R$ with $R=1 /(1-3 g R)$, namely

$$
R=\frac{1-\sqrt{1-12 g}}{6 g}
$$

$R$ is the gen. func. of quadrangulations with an origin and a marked edge

$$
\left.R\right|_{g^{n}}=3^{n} \operatorname{cat}(n)
$$

with

$$
\begin{gathered}
\operatorname{cat}(n) \equiv \frac{1}{n+1}\binom{2 n}{n} \\
\vec{Q}(n)=\frac{2}{n+2} \times 3^{n} \operatorname{cat}(n) \\
Q^{\bullet}(n)=\frac{1}{2 n} \times 3^{n} \operatorname{cat}(n) \\
Q(n)=\frac{1}{2 n(n+2)} \times 3^{n} \operatorname{cat}(n)
\end{gathered}
$$

## solution

$$
R_{\ell}=R \frac{\left(1-x^{\ell}\right)\left(1-x^{\ell+3}\right)}{\left(1-x^{\ell+1}\right)\left(1-x^{\ell+2}\right)}=R \frac{[\ell][\ell+3]}{[\ell+1][\ell+2]}
$$

where

$$
[\ell] \equiv \frac{1-x^{\ell}}{1-x}
$$

and where $x+x^{-1}+1=1 /\left(g R^{2}\right)$, namely

$$
x=\frac{1-24 g-\sqrt{1-12 g}+\sqrt{6} \sqrt{72 g^{2}+6 g+\sqrt{1-12 g}-1}}{2(6 g+\sqrt{1-12 g}-1)}
$$

## statistics of the distance between two points

## two-point function

a marked origin + a marked vertex at distance $m=d_{12}$ $\Leftrightarrow$ well-labeled tree with a marked vertex with label $m$

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$\diamond$ marked corner with label $m: R_{m}$

## two-point function

a marked origin + a marked vertex at distance $m=d_{12}$ $\Leftrightarrow$ well-labeled tree with a marked vertex with label $m$

$\diamond$ marked vertex with label $m: L_{m}=\log R_{m}$

## two-point function

a marked origin + a marked vertex at distance $m=d_{12}$ $\Leftrightarrow$ well-labeled tree with a marked vertex with label $m$

$\diamond$ marked vertex with label $m: L_{m}=\log R_{m}$
$\diamond$ impose $\min _{v \in \text { tree }} \ell(v) \geq 1$

## two-point function

a marked origin + a marked vertex at distance $m=d_{12}$ $\Leftrightarrow$ well-labeled tree with a marked vertex with label $m$

$\diamond$ marked vertex with label $m: L_{m}-L_{m-1}=\log \left(R_{m} / R_{m-1}\right)$
$\diamond$ impose $\min _{v \in \text { tree }} \ell(v)=1$

## two-point function

a marked origin + a marked vertex at distance $m=d_{12}$ $\Leftrightarrow$ well-labeled tree with a marked vertex with label $m$

$$
Q_{d_{12}}(g)=\left\{\begin{array}{cc}
\log \left(\frac{\left(\left[d_{12}\right]\right)^{2}\left[d_{12}+3\right]}{\left[d_{12}-1\right]\left(\left[d_{12}+2\right]\right)^{2}}\right) & \text { for } d_{12} \geq 2 \\
\log \left(R \frac{[1][4]}{[2][3]}\right) & \text { for } d_{12}=1
\end{array}\right.
$$

$\equiv$ generating function for doubly-pointed quadrangulations whose two marked (and distinguished) vertices are at distance $d_{12}$ from each other

## distance profile


$n=50$

## distance profile



$$
n=100
$$

distance profile

$n=150$

## distance profile


$n=200$

## distance profile


rescaled profiles

## local limit

## $<\mathrm{v}_{1}>(n)$


immediate neighbors

## local limit

$$
\begin{aligned}
& \left\langle\mathrm{V}_{2}\right\rangle(n) \\
& 10.8 \text {. }
\end{aligned}
$$

next-nearest neighbors

## local limit


next-next-nearest neighbors
limit laws
for large maps

## local limit

write

$$
\begin{gathered}
g=\frac{1}{12}\left(1-\epsilon^{2}\right) \\
R_{\ell}=\alpha_{\ell}+\beta_{\ell} \epsilon+\gamma_{\ell} \epsilon^{2}+\delta_{\ell} \epsilon^{3}+\cdots \\
\alpha_{\ell}=\frac{2 \ell(\ell+3)}{(\ell+1)(\ell+2)} \quad \beta_{\ell}=0 \quad \gamma_{\ell}=-\frac{\ell(\ell+3)\left(3 \ell^{2}+9 \ell-2\right)}{5(\ell+1)(\ell+2)} \\
\delta_{\ell}=\frac{\ell(\ell+3)\left(5 \ell^{4}+30 \ell^{3}+59 \ell^{2}+42 \ell+4\right)}{35(\ell+1)(\ell+2)}
\end{gathered}
$$

and the leading singularity (odd power in $\epsilon$ ) gives

$$
\left.R_{\ell}\right|_{g^{n}} \sim \frac{12^{n}}{\sqrt{\pi} n^{5 / 2}} \frac{3}{4} \delta_{\ell}
$$

$\ln [1]:=g=\frac{1}{12}\left(1-\epsilon^{2}\right) ; R=\frac{2}{1+\epsilon} ; x=X / . \operatorname{Sol} \bar{e}\left[X+\frac{1}{X}+1==\frac{1}{G R^{2}}, X\right][[1]]:$
$R 1:=R \frac{\left(1-x^{1}\right)\left(1-x^{1+3}\right)}{\left(1-x^{1+1}\right)\left(1-x^{1+2}\right)}$ :
Simplify[Series[R1, $\{\in, 0,3\}]]$
Out $[2]=\frac{21(3+1)}{2+31+1^{2}}-\frac{\left(1\left(-6+251+181^{2}+31^{3}\right)\right) \epsilon^{2}}{5\left(2+31+1^{2}\right)}+\frac{1\left(12+1301+2191^{2}+1491^{3}+451^{4}+51^{5}\right) \epsilon^{3}}{35\left(2+31+1^{2}\right)}+0[\epsilon]^{4}$
$\ln [3]:=$ Factor[CoefficientList [Hormal[x], e]]
Out $[3]=\left\{\frac{21(3+1)}{(1+1)(2+1)}, 0,-\frac{1(3+1)\left(-2+91+31^{2}\right)}{5(1+1)(2+1)}, \frac{1(3+1)\left(4+421+591^{2}+301^{3}+51^{4}\right)}{35(1+1)(2+1)}\right\}$
$\ln [4]:=\delta\left[1 \_\right]:=\frac{1(3+1)\left(4+421+591^{2}+301^{3}+51^{4}\right)}{35(1+1)(2+1)} ; \operatorname{Factor}\left[\frac{3}{2}(\delta[1]-\delta[1-1])\right]$
Out $[4]=\frac{6\left(-1+21+1^{2}\right)\left(4+141+271^{2}+201^{3}+51^{4}\right)}{351(1+1)(2+1)}$

## distance statistics

the average number $\left\langle e_{\ell}\right\rangle$ of edges at distance $\ell$ (i.e. $\ell-1 \leftrightarrow \ell$ ) in infinite quadrangulations is

$$
\left\langle e_{\ell}\right\rangle=\lim _{n \rightarrow \infty} \frac{\left.\left(R_{\ell}-R_{\ell-1}\right)\right|_{g^{n}}}{\left.R\right|_{g^{n}} /(2 n)}=\frac{3}{2}\left(\delta_{\ell}-\delta_{\ell-1}\right)
$$

one gets

$$
\begin{aligned}
\left\langle e_{\ell}\right\rangle=\frac{6}{35} \frac{\left(\ell^{2}+2 \ell-1\right)\left(5 \ell^{4}+20 \ell^{3}+27 \ell^{2}+\right.}{\ell(\ell+1)(\ell+2)} & \begin{array}{l}
\ell \ell+4) \\
\sim
\end{array} \frac{6}{7} \ell^{3}
\end{aligned}
$$

$\rightarrow$ fractal dimension $d_{F}=4$
NB: $\left\langle e_{1}\right\rangle=4$ obvious from Euler's relation

$$
\begin{gathered}
\log \left(R_{\ell}\right)=\tilde{\alpha}_{\ell}+\tilde{\beta}_{\ell} \epsilon+\tilde{\gamma}_{\ell} \epsilon^{2}+\tilde{\delta}_{\ell} \epsilon^{3}+\cdots \\
\tilde{\beta}_{\ell}=0 \quad \tilde{\delta}_{\ell}=\frac{5 \ell^{4}+30 \ell^{3}+59 \ell^{2}+42 \ell+4}{70}
\end{gathered}
$$

and the leading singularity gives

$$
\left.\log \left(R_{\ell}\right)\right|_{g^{n}} \sim \frac{12^{n}}{\sqrt{\pi} n^{5 / 2}} \frac{3}{4} \tilde{\delta}_{\ell}
$$

the average number $\left\langle v_{\ell}\right\rangle$ of vertices at distance $\ell$ in infinite quadrangulations is given by

$$
\begin{aligned}
\left\langle v_{\ell}\right\rangle=\frac{3}{35}\left((\ell+1)\left(5 \ell^{2}+10 \ell+2\right)\right. & \left.+\delta_{\ell, 1}\right) \\
& \stackrel{i \rightarrow \infty}{\sim} \frac{3}{7} \ell^{3}
\end{aligned}
$$

first values:

$$
\begin{array}{lll}
\left\langle e_{1}\right\rangle=4 & \left\langle e_{2}\right\rangle=19 & \left\langle e_{3}\right\rangle=\frac{1234}{25} \\
\left\langle v_{1}\right\rangle=3 & \left\langle v_{2}\right\rangle=\frac{54}{5} & \left\langle v_{3}\right\rangle=\frac{132}{5}
\end{array}
$$

## scaling limit

take $\ell$ large as $\ell=u \epsilon^{-1 / 2}$ with $u$ finite $\rightarrow$ scaling function $\mathcal{F}$ :

$$
R_{\ell}=2(1-\epsilon \mathcal{F}(u))+\mathcal{O}\left(\epsilon^{3 / 2}\right)
$$

whose small $u$ behavior can be read off the local limit

$$
\begin{gathered}
\alpha_{\ell}=2-\frac{4}{u^{2}} \epsilon+\mathcal{O}\left(\epsilon^{3 / 2}\right), \gamma_{\ell} \epsilon^{2}=-\frac{3 u^{2}}{5} \epsilon+\mathcal{O}\left(\epsilon^{3 / 2}\right), \delta_{\ell} \epsilon^{3}=\frac{u^{4}}{7} \epsilon+\mathcal{O}\left(\epsilon^{3 / 2}\right) \\
R_{\ell}=2-\epsilon\left(\frac{4}{u^{2}}+\frac{3 u^{2}}{5}-\frac{u^{4}}{7}+\mathcal{O}\left(u^{5}\right)\right)+\mathcal{O}\left(\epsilon^{3 / 2}\right)
\end{gathered}
$$

from the exact solution, one finds

$$
\mathcal{F}(u)=1+\frac{3}{\sinh ^{2}(\sqrt{3 / 2} u)}
$$

## scaling limit (fixed $n$ )

by a change of variables $g \rightarrow V \equiv g R$, we have

$$
R_{\ell} \left\lvert\, g^{n}=\oint \frac{d g}{2 i \pi g^{n+1}} R_{\ell}(g)=\oint \frac{d V(1-6 V)}{2 i \pi(V(1-3 V))^{n+1}} R_{\ell}(g)\right.
$$

for large $n$ and in the scaling limit

$$
\ell=r n^{1 / 4}
$$

do a saddle point calculation

$$
V=\frac{1}{6}\left(1+\mathrm{i} \frac{\xi}{\sqrt{n}}\right), \quad g=\frac{1}{12}\left(1+\frac{\xi^{2}}{n}\right)
$$

## use previous formulas with

$$
\begin{gathered}
\epsilon=\frac{-\mathrm{i} \xi}{\sqrt{n}} \quad u=r \sqrt{-\mathrm{i} \xi} \\
R_{\ell}=2\left(1+\frac{\mathrm{i} \xi}{\sqrt{n}} \mathcal{F}(r \sqrt{-\mathrm{i} \xi})\right) \quad \text { where } \mathcal{F}(u)=1+\frac{3}{\sinh ^{2}(\sqrt{3 / 2} u)}
\end{gathered}
$$

hence the large $n$ limit

$$
\left.R_{\ell}\right|_{g^{n}} \sim 2 \frac{12^{n}}{\pi n^{3 / 2}} \int_{-\infty}^{\infty} d \xi \xi^{2} e^{-\xi^{2}}(\mathcal{F}(r \sqrt{-\mathrm{i} \xi}))
$$

probability $\Phi(r)$ for a point (vertex or edge) to be at geodesic distance less than $r$ :

$$
\Phi(r)=\frac{4}{\sqrt{\pi}} \int_{0}^{\infty} d \xi \xi^{2} e^{-\xi^{2}}\left\{1-6 \frac{1-\cosh (r \sqrt{3 \xi}) \cos (r \sqrt{3 \xi})}{(\cosh (r \sqrt{3 \xi})-\cos (r \sqrt{3 \xi}))^{2}}\right\}
$$

probability density $\rho(r)$ for a point (vertex or edge) to be at geodesic distance $r$

$$
\rho(r)=\frac{d \Phi(r)}{d r}
$$

$\Phi(r)$

$\rho(r)$


$$
\rho(r) \stackrel{r \rightarrow 0}{\sim} \frac{3}{7} r^{3}, \quad \rho(r) \stackrel{r \rightarrow \infty}{\sim} \exp \left(-\frac{3}{4} 3^{2 / 3} r^{4 / 3}\right)
$$

in agreement with $\left\langle v_{\ell}\right\rangle$ and Fisher's law $\delta=\frac{4}{3}=\frac{1}{1-\nu}$ with $\nu=\frac{1}{4}=\frac{1}{d_{F}}$
statistics of geodesics

## quadrang. with a marked geodesic








$$
U_{2}(g)=g+10 g^{2}+\ldots
$$




## collection of geodesics



## local limit

$$
\begin{gathered}
Z_{\ell}=A_{\ell}-C_{\ell} \epsilon^{2}+\frac{2}{3} D_{\ell} \epsilon^{3}+\cdots \\
A_{\ell}=2^{\ell} \frac{l+3}{3(\ell+1)}, \quad D_{\ell}=2^{\ell} \frac{\ell(\ell+2)(\ell+3)(\ell+4)\left(3 \ell^{2}+12 \ell+13\right)}{420(\ell+1)} \\
U_{\ell}=a_{\ell}-c_{\ell} \epsilon^{2}+\frac{2}{3} d_{\ell} \epsilon^{3}+\cdots \\
a_{\ell}=A_{\ell}-\sum_{j=1}^{\ell-1} a_{j} A_{\ell-j}, \quad d_{\ell}=D_{\ell}-\sum_{j=1}^{\ell-1}\left(a_{j} D_{\ell-j}+d_{j} A_{\ell-j}\right)
\end{gathered}
$$

introduce $\hat{A}(t)=\sum_{\ell \geq 1} A_{\ell} t^{\ell}, \ldots$

$$
\hat{a}(t)=\frac{\hat{A}(t)}{1+\hat{A}(t)}, \quad \hat{d}(t)=\frac{\hat{D}(t)}{(1+\hat{A}(t))^{2}}
$$

## $\rightarrow$ exact expression for $\hat{d}(t)$

$$
\begin{aligned}
\hat{d}(t)= & 4 t+\frac{80}{3} t^{2}+132 t^{3}+\cdots \\
& \langle\text { geods }\rangle_{1}=d_{1}=4 \\
& \langle\text { geods }\rangle_{2}=d_{2}=\frac{80}{3} \\
& \langle\text { geods }\rangle_{3}=d_{3}=132
\end{aligned}
$$

large $\ell$ behavior of $d_{\ell}$ ?

$$
\begin{aligned}
& A_{\ell} \sim \frac{2^{\ell}}{3}\left(1+\frac{2}{\ell}\right) \rightarrow \hat{A}(t) \sim \frac{1}{3(1-2 t)}-\frac{2}{3} \log (1-2 t) \\
& D_{\ell} \sim \frac{2^{\ell} \ell^{5}}{140} \rightarrow \hat{D}(t) \sim \frac{6}{7(1-2 t)^{6}} \\
& \hat{a}(t) \sim 1-3(1-2 t)-6(1-2 t)^{2} \log (1-2 t) \rightarrow a_{\ell} \sim \frac{2^{\ell} 12}{\ell^{3}} \\
& \hat{d}(t) \sim \frac{54}{7(1-2 t)^{4}} \rightarrow d_{\ell} \sim \frac{2^{\ell} 9 \ell^{3}}{7} \\
& d_{\ell} \\
& \sim\left(3 \times 2^{\ell}\right) \times \frac{3}{7} \ell^{3}
\end{aligned}
$$

## collection of geodesics

$$
\begin{gathered}
U_{\ell}^{(k)}=a_{\ell}^{(k)}-c_{\ell}^{(k)} \epsilon^{2}+\frac{2}{3} d_{\ell}^{(k)} \epsilon^{3}+\cdots \\
d_{\ell}^{(k)}=k \times\left(3 \times 2^{\ell}\right)^{k} \times \frac{3}{7} \ell^{3} \\
\tilde{U}_{\ell}^{(k)}=\tilde{a}_{\ell}^{(k)}-\tilde{c}_{\ell}^{(k)} \epsilon^{2}+\frac{2}{3} \tilde{d}_{\ell}^{(k)} \epsilon^{3}+\cdots \\
\tilde{d}_{\ell}^{(k)}=k\left(a_{\ell}\right)^{k-1} d_{\ell} \\
\tilde{d}_{\ell}^{(k)}=k \times\left(3 \times 2^{\ell}\right)^{k} \times 4^{k-1} \frac{3}{7} \ell^{6-3 k}
\end{gathered}
$$

number of vertices at distance $\ell$ reached by $k$ avoiding geods $=4^{k-1} \frac{3}{7} \ell^{6-3 k}$

## scaling limit (exponents)

the average number of pairs of points linked by $k$ avoiding geods and at rescaled distance in the range $[r, r+d r]$ behaves as

$$
n \times\left(n^{1 / 4}\right)^{6-3 k} n^{1 / 4} d r \times \rho^{(k)}(r)
$$

with

$$
\begin{gathered}
\rho^{(k)}(r)^{r \rightarrow 0}{ }_{\sim}^{\sim} r^{6-3 k} \\
\Rightarrow n^{(11-3 k) / 4}
\end{gathered}
$$

$k=1: \quad n^{2}$
$k=2: \quad n^{5 / 4}$
$k=3: n^{1 / 2}$

## scaling limit (scaling functions)

$$
\begin{aligned}
& g=\frac{1}{12}\left(1-\epsilon^{2}\right), \quad \ell=u \epsilon^{-1 / 2} \\
& R_{\ell} \sim 2(1-\epsilon \mathcal{F}(u)), \quad \frac{Z_{\ell}}{2^{\ell}} \sim \frac{1}{3}+\epsilon^{1 / 2} \mathcal{H}(u), \quad \frac{U_{\ell}}{2^{\ell}} \sim \epsilon^{3 / 2} \mathcal{L}(u) \\
& \quad \mathcal{F}(u)=-3 \frac{d}{d u} \mathcal{H}(u), \quad \mathcal{L}(u)=9 \frac{d^{2}}{d u^{2}} \mathcal{H}(u)=-3 \frac{d}{d u} \mathcal{F}(u)
\end{aligned}
$$

scaling limit $\ell=r n^{1 / 4}$

$$
\begin{aligned}
\left.U_{\ell}\right|_{g^{n}} & \sim \frac{12^{n}}{2 \sqrt{\pi} n^{5 / 2}}\left(3 \cdot 2^{\ell}\right) n \rho(r) \frac{1}{n^{1 / 4}} \\
\left.U_{\ell}^{(k)}\right|_{g^{n}} & \sim \frac{12^{n}}{2 \sqrt{\pi} n^{5 / 2}} k\left(3 \cdot 2^{\ell}\right)^{k} n \rho(r) \frac{1}{n^{1 / 4}}
\end{aligned}
$$

$\rightarrow$ no new scaling function

$$
\begin{gathered}
\frac{\tilde{U}_{\ell}^{(2)}}{2^{2 \ell}}=\left(\frac{U_{\ell}}{2^{\ell}}\right)^{2} \sim \epsilon^{3}(\mathcal{L}(u))^{2} \\
n \tilde{U}_{\ell}^{(2)}{\mid g^{n}}^{\sim} \frac{12^{n}}{2 \sqrt{\pi} n^{5 / 2}} 2\left(3 \cdot 2^{\ell}\right)^{2} n^{5 / 4} \rho^{(2)}(r) \frac{1}{n^{1 / 4}}
\end{gathered}
$$

with the new scaling function

$$
\rho^{(2)}(r)=\frac{1}{9 \sqrt{\pi}} \int_{-\infty}^{\infty} d \xi \xi^{4} e^{-\xi^{2}}(\mathcal{L}(r \sqrt{-\mathrm{i} \xi}))^{2}
$$

$$
\rho^{(2)}(r)
$$



$$
\begin{gathered}
\frac{\tilde{U}_{\ell}^{(3)}}{2^{3 \ell}}=\left(\frac{U_{\ell}}{2^{\ell}}\right)^{3} \sim \epsilon^{9 / 2}(\mathcal{L}(u))^{3} \\
n \tilde{U}_{\ell}^{(3)}{\mid g^{n}}^{\sim} \frac{12^{n}}{2 \sqrt{\pi} n^{5 / 2}} 3\left(3 \cdot 2^{\ell}\right)^{3} n^{1 / 2} \rho^{(3)}(r) \frac{1}{n^{1 / 4}}
\end{gathered}
$$

with the new scaling function

$$
\rho^{(3)}(r)=\frac{2}{81 \sqrt{\pi}} \int_{-\infty}^{\infty} d \xi \frac{\xi^{5}}{\mathrm{i}} e^{-\xi^{2}} \sqrt{-\mathrm{i} \xi}(\mathcal{L}(r \sqrt{-\mathrm{i} \xi}))^{3}
$$

$$
\rho^{(3)}(r)
$$



$$
\rho^{(3)}(r)
$$

## three-point statistics

## Miermont's bijection

start with a multiply-pointed planar quadrangulation with $p$ marked vertices (=sources) distinguished as $v_{1}, \cdots, v_{p}$ and satisfying $d\left(v_{i}, v_{j}\right) \geq 2$


## Miermont's bijection

natural labeling: $\ell(v) \equiv \min _{j=1, \ldots . p} d\left(v, v_{j}\right)$

one can favor/penalize some of the sources by attaching to each source $v_{i}$ a delay $\tau_{i}=$ integer

this defines a "delayed distance" from $v$ to the source $v_{j}$ :

$$
\ell_{j}(v) \equiv d\left(v, v_{j}\right)+\tau_{j}
$$

a vertex $v$ now receives the label:

$$
\ell(v) \equiv \min _{j=1, \ldots p} \ell_{j}(v)=\min _{j=1, \ldots p}\left(d\left(v, v_{j}\right)+\tau_{j}\right)
$$

which is the "distance" to the closest source, where the distance from $v_{j}$ incorporates a penalty $\tau_{j}$
we choose delays so that:

$$
\diamond\left|\tau_{i}-\tau_{j}\right|<d\left(v_{i}, v_{j}\right) \forall i \neq j \quad \text { (cond. 1) }
$$

$\rightarrow$ ensures that $\ell\left(v_{i}\right)=\tau_{i}$

$$
\diamond \tau_{i}-\tau_{j}=d\left(v_{i}, v_{j}\right) \bmod 2 \quad \text { (cond. 2) }
$$

$\rightarrow$ ensures that the parity of $\ell_{j}(v)$ is independent of $j$ so that again $\left|\ell(v)-\ell\left(v^{\prime}\right)\right|=1$ for $v$ and $v^{\prime}$ neighbors

faces $\rightarrow$ edges


## faces $\rightarrow$ edges


same rules as in Schaeffer's bijection



end up with a planar well-labeled map with $p$ faces

## well-labeled maps



## well-labeled maps


$\diamond$ labels vary by at most 1 between neighbors
$\left|\ell(v)-\ell\left(v^{\prime}\right)\right| \leq 1$ if $v$ and $v^{\prime}$ are neighbors on the map

## well-labeled maps


$\diamond$ labels vary by at most 1 between neighbors
$\left|\ell(v)-\ell\left(v^{\prime}\right)\right| \leq 1$ if $v$ and $v^{\prime}$ are neighbors on the map
$\diamond \min _{v \text { incident to } F_{i}} \ell(v)=1+\tau_{i}$
bijection: for fixed given delays
$p$-pointed quadrangulations
with marked vertices satisfying
$\diamond d\left(v_{i}, v_{j}\right)>\left|\tau_{i}-\tau_{j}\right| \forall i \neq j$
$\diamond d\left(v_{i}, v_{j}\right)=\tau_{i}-\tau_{j} \bmod 2$
well-labeled maps with $p$ faces
with labels satisfying
$\diamond\left|\ell(v)-\ell\left(v^{\prime}\right)\right| \leq 1$ if $v$ and $v^{\prime}$ are neighbors
$\diamond \min _{v \text { incident to } F_{i}} \ell(v)=1+\tau_{i}$
this coding keeps track of some of the distances:
if $v$ is incident to $F_{i}$, then the minimum of $\ell_{j}=d\left(v, v_{j}\right)+\tau_{j}$ is reached for $j=i$ and therefore:

$$
d\left(v, v_{i}\right)=\ell(v)-\tau_{i}
$$


planar maps are classified according to their backbone

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planar maps are classified according to their backbone

all vertices have degree $\geq 3 \Rightarrow$ finite number of backbones

## case of 3 marked vertices

planar maps with 3 (distinguished) faces
$\rightarrow$ seven possible backbones

map $=$ backbone + attached trees


## distance parametrization

for 3 points, we can use the following parametrization:

$$
\begin{aligned}
& d_{12} \equiv d\left(v_{1}, v_{2}\right)=s+t \\
& d_{23} \equiv d\left(v_{2}, v_{3}\right)=t+u \\
& d_{31} \equiv d\left(v_{3}, v_{1}\right)=u+s
\end{aligned}
$$


with $s, t, u$ integers, $s, t, u \geq 0$ and at most one may vanish

## choice of delays

idea: relate the delays to the distances, namely choose:

$$
\tau_{1}=-s, \quad \tau_{2}=-t, \quad \tau_{3}=-u
$$

$\diamond \tau_{1}-\tau_{2}=-s+t=s+t \bmod 2=d\left(v_{1}, v_{2}\right) \bmod 2$
$\diamond\left|\tau_{1}-\tau_{2}\right|=\left|d_{23}-d_{31}\right| \leq d_{12}$ (triangular inequalities) and equality only if the 3 vertices are "aligned": for instance $d_{23}-d_{31}=d_{12}$ only if $v_{1}$ lies on a geodesic path between $v_{2}$ and $v_{3}$ (i.e. $s=0$ )

- assume that the 3 vertices are not aligned $\Leftrightarrow s, t, u>0$
- treat later the case of aligned vertices (includes the case when two vertices are immediate neighbors)
new constraint on labels:
$\diamond$ vertex on the boundary between two faces

length $=\ell(v)-\tau_{1}+\ell(v)-\tau_{2}=2 \ell(v)+(s+t) \geq s+t=d_{12}$
$\Rightarrow \ell(v) \geq 0$ for vertices on boundaries of the well-labeled map


$$
\begin{aligned}
s+t=d_{12} & =\ell(v)-\tau_{1}+\ell\left(v^{\prime}\right)-\tau_{2}+d\left(v, v^{\prime}\right) \\
& =\ell(v)+\ell\left(v^{\prime}\right)+d\left(v, v^{\prime}\right)+s+t
\end{aligned}
$$

$\Rightarrow \ell(v)+\ell\left(v^{\prime}\right)+d\left(v, v^{\prime}\right)=0$
$\Rightarrow v=v^{\prime}$ and $\ell(v)=0$
$F_{1}$ and $F_{2}$ must have a common boundary (+ permutations) $\exists$ label 0 on each boundary
rules out backbones of type:

the only possible backbones are those of type

and those of type

in practice, the latter can be viewed as degenerate cases of the former when one of the boundaries reduces to a single vertex.

## rules on labels


bijection:
triply-pointed quadrangulations
with marked vertices at prescribed pairwise distances $d_{12}=s+t, d_{23}=t+u$ and $d_{31}=u+s$ with $s, t, u>0$

$$
\Uparrow
$$

well-labeled maps with 3 faces
with a backbone

or its degenerate versions

bijection:
triply-pointed quadrangulations
with marked vertices at prescribed pairwise distances $d_{12}=s+t, d_{23}=t+u$ and $d_{31}=u+s$

$$
\Uparrow
$$

well-labeled maps with 3 faces
with a backbone

or its degenerate versions
or its degenerate versions


## aligned case

if $v_{3}$ lies between $v_{1}$ and $v_{2} \quad\left(d_{31}=s, d_{23}=t, d_{12}=s+t\right)$ apply Miermont's construction on $v_{1}$ and $v_{2}$ only, with delays $\tau_{1}=-s$ and $\tau_{2}=-t$


## enumeration of well-labeled maps


reminder: the generating function for well-labeled trees planted at a label $\ell$ and with the condition:

$$
\min _{v \in \text { tree }} \ell(v) \geq 1
$$

is given by

$$
R_{\ell}=R \frac{[\ell][\ell+3]}{[\ell+1][\ell+2]} \quad \text { where }[\ell] \equiv \frac{1-x^{\ell}}{1-x}
$$

reminder: the generating function for well-labeled trees planted at a label $\ell$ and with the condition:

$$
\min _{v \in \text { tree }} \ell(v) \geq 1
$$

is given by

$$
R_{\ell}=R \frac{[\ell][\ell+3]}{[\ell+1][\ell+2]} \quad \text { where }[\ell] \equiv \frac{1-x^{\ell}}{1-x}
$$

if we wish instead: $\min _{v \in \text { tree }} \ell(v) \geq 1-s$
this generating function is nothing but: $R_{\ell+s}$
as obtained by a simple shift by $s$ of all labels
consider the generating function $F_{s, t, u}(g)$

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consider the generating function $F_{s, t, u}(g)$
first 0


$$
X_{s, t} \quad X_{t, u} \quad X_{u, s}
$$


consider the generating function $F_{s, t, u}(g)$

$$
Y_{s, t, u}
$$

$$
X_{S, t}
$$

$$
X_{t, u}
$$

$$
X_{u, s}
$$

$$
Y_{s, t, u}
$$

$$
F_{s, t, u}(g)=X_{s, t} X_{t, u} X_{u, s}\left(Y_{s, t, u}\right)^{2}
$$

## $Y_{s, t, u}$

$X_{S, t}$

## $X_{t, u}$

$X_{u, s}$
$Y_{s, t, u}$


$$
X_{s, t}=\sum_{\substack{\text { motzkin paths of length } m \\ \mathcal{M}=\left(0=\ell_{0}, \ell_{1}, \ldots, \ell_{m}=0\right)}} \prod_{\substack{k=0}} g R_{\ell_{k}+s} R_{\ell_{k}+t}
$$

## recursion relation:

$$
X_{s, t}=1+g R_{s} R_{t} X_{s, t}\left(1+g R_{s+1} R_{t+1} X_{s+1, t+1}\right)
$$

## recursion relation:

$$
X_{s, t}=1+g R_{s} R_{t} X_{s, t}\left(1+g R_{s+1} R_{t+1} X_{s+1, t+1}\right)
$$

solution:

$$
X_{s, t}=\frac{[3][s+1][t+1][s+t+3]}{[1][s+3][t+3][s+t+1]}
$$

## recursion relation:

$$
X_{s, t}=1+g R_{s} R_{t} X_{s, t}\left(1+g R_{s+1} R_{t+1} X_{s+1, t+1}\right)
$$

solution:

$$
X_{s, t}=\frac{[3][s+1][t+1][s+t+3]}{[1][s+3][t+3][s+t+1]}
$$

similarly, recursion relation for $Y_{s, t, u}$ :
$Y_{s, t, u}=1+g^{3} R_{s} R_{t} R_{u} R_{s+1} R_{t+1} R_{u+1}$

$$
\times X_{s+1, t+1} X_{t+1, u+1} X_{u+1, s+1} Y_{s+1, t+1, u+1}
$$

## recursion relation:

$$
X_{s, t}=1+g R_{s} R_{t} X_{s, t}\left(1+g R_{s+1} R_{t+1} X_{s+1, t+1}\right)
$$

solution:

$$
X_{s, t}=\frac{[3][s+1][t+1][s+t+3]}{[1][s+3][t+3][s+t+1]}
$$

similarly, recursion relation for $Y_{s, t, u}$ :
$Y_{s, t, u}=1+g^{3} R_{s} R_{t} R_{u} R_{s+1} R_{t+1} R_{u+1}$

$$
\times X_{s+1, t+1} X_{t+1, u+1} X_{u+1, s+1} Y_{s+1, t+1, u+1}
$$

solution:

$$
Y_{s, t, u}=\frac{[s+3][t+3][u+3][s+t+u+3]}{[3][s+t+3][t+u+3][u+s+3]}
$$









$$
X_{s, t} Z_{s} Z_{t}=Z_{s+t}
$$

namely

$$
X_{s, t}=\frac{Z_{s+t}}{Z_{s} Z_{t}}
$$

with

$$
Z_{i}=\frac{[1][i+3]}{[3][i+1]}
$$

so that

$$
X_{s, t}=\frac{[3][s+1][t+1][s+t+3]}{[1][s+3][t+3][s+t+1]}
$$

## three-point function

$$
\begin{aligned}
& F_{s, t, u}(g)=X_{s, t} X_{t, u} X_{u, s}\left(Y_{s, t, u}\right)^{2} \\
& =\frac{[3]([s+1][t+1][u+1][s+t+u+3])^{2}}{[1]^{3}[s+t+1][s+t+3][t+u+1][t+u+3][u+s+1][u+s+3]}
\end{aligned}
$$

and the three-point function for quadrangulations reads

$$
Q_{d_{12}, d_{23}, d_{31}}(g)=\Delta_{s} \Delta_{t} \Delta_{u} F_{s, t, u}(g)
$$

with $\Delta_{s} f(s) \equiv f(s)-f(s-1)$, and
$s=\left(d_{12}-d_{23}+d_{31}\right) / 2$
$t=\left(d_{12}+d_{23}-d_{31}\right) / 2$
$u=\left(-d_{12}+d_{23}+d_{31}\right) / 2$

## scaling limit

$$
g=\frac{1}{12}\left(1-\epsilon^{2}\right), \quad \ell=L \epsilon^{-1 / 2} \quad \text { with } \epsilon \rightarrow 0
$$

replace in any well-balanced combination of [.]'s:

$$
\begin{aligned}
& {[\ell] }=\frac{1-x^{\ell}}{1-x} \rightarrow \sinh (\alpha L), \quad \alpha=\sqrt{\frac{3}{2}} \\
& s=S \epsilon^{-1 / 2}, t=T \epsilon^{-1 / 2}, u=U \epsilon^{-1 / 2} \\
& X_{s, t} \rightarrow 3, \quad Y_{s, t, u} \rightarrow \epsilon^{-1 / 2} \mathcal{Y}(S, T, U ; \sqrt{3 / 2})
\end{aligned}
$$

with

$$
\mathcal{Y}(S, T, U ; \alpha) \equiv \frac{1}{3 \alpha} \frac{\sinh \alpha S \sinh \alpha T \sinh \alpha U \sinh \alpha(S+T+U)}{\sinh \alpha(S+T) \sinh \alpha(T+U) \sinh \alpha(U+S)}
$$

$$
F_{s, t, u}(g) \sim \epsilon^{-1} \mathcal{F}(S, T, U ; \sqrt{3 / 2})
$$

with $\mathcal{F}(S, T, U ; \alpha)=$

$$
\begin{gathered}
\frac{3}{\alpha^{2}}\left(\frac{\sinh (\alpha(S+T+U)) \sinh (\alpha S) \sinh (\alpha T) \sinh (\alpha U)}{\sinh (\alpha(S+T)) \sinh (\alpha(T+U)) \sinh (\alpha(U+S))}\right)^{2} \\
Q_{d_{12}, d_{23}, d_{31}}(g) \sim \epsilon^{1 / 2} \mathcal{Q}\left(D_{12}, D_{23}, D_{31} ; \sqrt{3 / 2}\right)
\end{gathered}
$$

with

$$
\mathcal{Q}\left(D_{12}, D_{23}, D_{31} ; \alpha\right)=\frac{1}{2} \partial_{S} \partial_{T} \partial_{U} \mathcal{F}(S, T, U ; \alpha)
$$

$$
S=\left(D_{12}-D_{23}+D_{31}\right) / 2
$$

$$
T=\left(D_{12}+D_{23}-D_{31}\right) / 2
$$

$$
U=\left(-D_{12}+D_{23}+D_{31}\right) / 2
$$

## three-point function

$$
\begin{gathered}
d_{12}=D_{12} n^{1 / 4}
\end{gathered} d_{23}=D_{23} n^{1 / 4} \quad d_{31}=D_{31} n^{1 / 4}
$$

we get the probability density
$\rho\left(D_{12}, D_{23}, D_{31}\right)=\frac{2}{\sqrt{\pi}} \int_{-\infty}^{\infty} d \xi \frac{\xi}{\mathrm{i}} e^{-\xi^{2}} \mathcal{Q}\left(D_{12}, D_{23}, D_{31} ; \sqrt{3 / 2} \sqrt{-\mathrm{i} \xi}\right)$
$\rho\left(D_{12}, D_{23}, D_{31}\right) d D_{12} d D_{23} d D_{31}$ is the probability that the three marked vertices be at rescaled pairwise distances in the ranges $\left[D_{12}, D_{12}+d D_{12}\right],\left[D_{23}, D_{23}+d D_{23}\right]$, [ $\left.D_{31}, D_{31}+d D_{31}\right]$, in the ensemble of triply-pointed quadrangulations of fixed large size $n$

## conditional probability density

fix one of the distances, say $D_{12}$, and consider the conditional probability density

$$
\rho\left(D_{23}, D_{31} \mid D_{12}\right) \equiv \frac{\rho\left(D_{12}, D_{23}, D_{31}\right)}{\rho\left(D_{12}\right)}
$$

$\rho\left(D_{23}, D_{31} \mid D_{12}\right) d D_{23} d D_{31}$ is the probability that the third marked vertex be at rescaled distances in the ranges $\left[D_{23}, D_{23}+d D_{23}\right]$ and $\left[D_{31}, D_{31}+d D_{31}\right]$ from the first two marked vertices in the ensemble of triply-pointed quadrangulations of fixed large size $n$, given that the distance between the first two marked vertices is $D_{12}$

## $D_{12}=0.8$



## $D_{12}=1.5$



## $D_{12}=3.0$



## limit of small $D_{12}$



$$
\rho\left(D_{23}, D_{31} \mid D_{12}\right) \sim \frac{1}{D_{12}} \times \rho(D) \times \psi(\omega)
$$

where $D=\left(D_{23}+D_{31}\right) / 2, \omega=\left(D_{31}-D_{23}\right) / D_{12}$, and

$$
\psi(\omega) \equiv \frac{21}{64}\left(1-\omega^{2}\right)^{2}\left(3-\omega^{2}\right)
$$




## limit of large $D_{12}$



$$
\rho\left(D_{23}, D_{31} \mid D_{12}\right) \sim \frac{1}{2 D_{12}} \times\left(9 D_{12}\right)^{1 / 3} \varphi(\nu)
$$

where $\nu=\left(9 D_{12}\right)^{1 / 3}\left(D_{23}+D_{31}-D_{12}\right) / 2$, and

$$
\varphi(\nu) \equiv \frac{4}{3} \sinh (\nu / 2)^{2}\left(11 e^{-2 \nu}-8 e^{-3 \nu}\right)
$$



## other applications

phenomenon of confluence of geodesics (Le Gall)
$v_{3}$

## other applications

phenomenon of confluence of geodesics (Le Gall)

what is the law for the length $\delta$ of the common part ?
geodesics triangle

$v_{1}$
geodesics triangle

$v_{1}$
joint law for the 6 lengths and 2 areas ?


$$
X_{s^{\prime \prime}, t^{\prime \prime}} X_{t^{\prime \prime}, u^{\prime \prime}} X_{u^{\prime \prime}, s^{\prime \prime}} Y_{s^{\prime}, t^{\prime}, u^{\prime}} Y_{s^{\prime \prime}, t^{\prime \prime}, u^{\prime \prime}}
$$



law for the length $\delta$ of the common part

law for the length $\delta$ of the common part
$\rho(\delta)$

$\delta$
law for the length $\Delta_{12}$ of the open part from $v_{1}$ to $v_{2}$

$v_{1}$
law for the length $\Delta_{12}$ of the open part from $v_{1}$ to $v_{2}$

joint law for the lengths $\delta_{1}$ and $\delta_{2}$ given $D_{12}$

$v_{1}$
joint law for the lengths $\delta_{1}$ and $\delta_{2}$ given $D_{12}$

joint law for the lengths $\delta_{1}$ and $\delta_{2}$ given $D_{12}$

joint law for the lengths $\delta_{1}$ and $\delta_{2}$ given $D_{12}$



$\Delta_{s^{\prime}} \Delta_{t^{\prime}} \Delta_{u^{\prime}} Y_{s^{\prime}, t^{\prime}, u^{\prime}} \rightarrow$ triangle with geodesic boundaries of side lengths $s^{\prime}+t^{\prime}, t^{\prime}+u^{\prime}, u^{\prime}+s^{\prime}$

$\partial_{S^{\prime}} \partial_{T^{\prime}} \partial_{U^{\prime}} \mathcal{Y}\left(S^{\prime}, T^{\prime}, U^{\prime}\right) \rightarrow$ triangle with geodesic boundaries of side lengths $S^{\prime}+T^{\prime}, T^{\prime}+U^{\prime}, U^{\prime}+S^{\prime}$

another triangle with geodesic boundaries of different side lengths

make the side lengths of the two triangles identical (here $S^{\prime \prime}>S^{\prime}, T^{\prime \prime}>U^{\prime}$ and $U^{\prime \prime}>U^{\prime}$ )


$\rightarrow$ one contribution to the three-point function


## a word on branching processes

## a branching process

## a random map is the "superposition" of

## a branching process

a random map is the "superposition" of

- a random planar tree


## a branching process

a random map is the "superposition" of

- a random planar tree
- integer labels on the tree


## a branching process

a random map is the "superposition" of

- a random planar tree
- integer labels on the tree
- boundary condition (positive labels)


## a branching process

a random map is the "superposition" of

- a random planar tree $\rightarrow$ genealogical tree
- integer labels on the tree
- boundary condition (positive labels)
a parent individual gives rise to $k$ children with probability $p(k)=(1-p) p^{k}, \quad$ (average number of children $\frac{p}{1-p}$ )


## a branching process

a random map is the "superposition" of

- a random planar tree $\rightarrow$ genealogical tree
- integer labels on the tree $\rightarrow$ diffusion process in 1D
- boundary condition (positive labels)
a parent individual gives rise to $k$ children with probability $p(k)=(1-p) p^{k}, \quad$ (average number of children $\frac{p}{1-p}$ )
the child of a parent at position $\ell$ lives at position $\ell, \ell \pm 1$


## a branching process

a random map is the "superposition" of

- a random planar tree $\rightarrow$ genealogical tree
- integer labels on the tree $\rightarrow$ diffusion process in 1D
- boundary condition (positive labels) $\rightarrow$ walls, forbidden zone
a parent individual gives rise to $k$ children with probability $p(k)=(1-p) p^{k}, \quad$ (average number of children $\frac{p}{1-p}$ )
the child of a parent at position $\ell$ lives at position $\ell, \ell \pm 1$
what is the probability $\mathcal{P}_{\ell}(p)$ for the population whose germ is at position $\ell$ to reach position 0 ?
$\mathcal{P}_{\ell}(p)=1-(1-p) R_{\ell}(g)$ with $g=\frac{p(1-p)}{3}$

$$
\mathcal{P}_{\ell}(p)=1-\frac{1-|2 p-1|}{2 p} \frac{\left(1-x^{\ell}\right)\left(1-x^{\ell+3}\right)}{\left(1-x^{\ell+1}\right)\left(1-x^{\ell+2}\right)}
$$

with $x=\frac{1+2|1-2 p|-\sqrt{3|1-2 p|} \sqrt{2+|1-2 p|}}{1-|1-2 p|}$
$\mathcal{P}_{\ell}(p) \stackrel{\ell \rightarrow \infty}{\sim} S(p)$ : survival probability


$$
S(p)=1-\frac{1-|2 p-1|}{2 p}=\left\{\begin{array}{cl}
0 & p \leq \frac{1}{2} \\
\frac{2 p-1}{p} & p \geq \frac{1}{2}
\end{array}\right.
$$



scaling behavior around $p=\frac{1}{2}$ :

$$
\mathcal{P}_{\ell}(p) \sim|2 p-1|\left(\frac{3}{\sinh ^{2}\left(\sqrt{3 / 2} \ell|2 p-1|^{1 / 2}\right)}+1\right)+(2 p-1)
$$

## THE END

